

Review

Friday, September 30, 2022 4:17 PM

Quiz B1

How many comparisons does the insertion sort need to sort this array?

$A = [3, 5, 1, 8, 2, 6]$

<u>i</u>	<u>a</u>		<u>C</u>	<u>A</u>
2	5	$\begin{array}{cccccc} & \downarrow & \downarrow \\ \underline{3} & 5 & 1 & 8 & 2 & 6 \\ \underline{3} & 5 & & & & \end{array}$	1	2
3	1	$\begin{array}{cccccc} & & & & & \\ & 1 & 3 & 5 & & \end{array}$	2	4
4	8	$\begin{array}{cccccc} & & & & & \\ & 1 & 3 & 5 & 8 & \end{array}$	1	2
5	2	$\begin{array}{cccccc} & & & & & \\ & 1 & 2 & 3 & 5 & 8 \end{array}$	4	5
6	6	$\begin{array}{cccccc} & & & & & \\ & 1 & 2 & 3 & 5 & 6 & 8 \end{array}$	2	3
			<hr/> 10	<hr/> 16

How many comparisons does the Bottom-up Mergesort need to sort this array

$A = [2, 4, 3, 5, 6, 7]$

<u>2 4 3 5 6 7</u>	<u>C</u>	<u>A</u>
$\begin{array}{ccccc} \underline{2} & \underline{4} & \underline{3} & \underline{5} & & \underline{6} & \underline{7} \end{array}$	3	12
$\begin{array}{ccccc} \underline{2} & \underline{3} & \underline{4} & \underline{5} & & & \end{array}$	3	8
$\begin{array}{ccccc} \underline{2} & \underline{3} & \underline{4} & \underline{5} & \underline{6} & \underline{7} \end{array}$	4	12
	<hr/> 10	<hr/> 32

Let $T(n) = 5T(n/6) + n / \log n$.

Apply the master theorem to give an asymptotic bound for $T(n)$ using an appropriate notation (O , Ω , or Θ).

$$R = \log_6 5 \approx 0.9$$

$$\frac{n}{\log n} = \Omega(n^{R+\epsilon})$$

\Rightarrow Case :

$$T(n) = \Theta(n / \log n)$$

HW 2

Exer 29.

$$n! = \Theta(n^n) \quad \text{False} \quad n! = o(n^n)$$

$$n! = 1 \cdot 2 \cdot 3 \cdot \dots \cdot n$$

$$< 1 \cdot 2 \cdot 3 \cdot n \cdot n \dots n$$

$$= 6 \cdot n^{n-3} = o(n^3 \cdot n^{n-3}) = o(n^n)$$

Note: $\log(n!) = \Theta(\log n^n) = \Theta(n \log n)$

1.15. Express the following functions in terms of the Θ -notation.

(a) $\underline{2n} + \underline{3 \log^{100} n} = \Theta(n)$

(b) $\underline{7n^3} + \underline{1000n \log n} + \underline{3n} = \Theta(n^3)$

(c) $\underline{3n^{1.5}} + (\sqrt{n})^3 \log n = \Theta(n^{1.5} \log n)$

(d) $\underline{2^n} + \underline{100^n} + \underline{n!} = \Theta(n!)$

1.16. Express the following functions in terms of the Θ -notation.

(a) $18n^3 + \log n^8 = \Theta(n^3)$

(b) $(n^3 + n)/(n + 5) = n^2 + \dots = \Theta(n^2)$

$(\log n)^2$

(c) $\underline{\log^2 n} + \underline{\sqrt{n}} + \underline{\log \log n} = \Theta(\sqrt{n})$

(d) $\underline{n!/2^n} + \underline{n^n} = \Theta(n^n)$

(e) $n! + n^n/2^n = \Theta\left(\left(\frac{n}{2}\right)^n\right)$

$n! \sim n^n/2^n$
 $n=5 \rightarrow 120$
 $n=6 \rightarrow 720$
 $\frac{6^6}{64} = 729$

1.22. Show that $n^{100} = O(2^n)$, but $2^n \neq O(n^{100})$.

1.23. Show that 2^n is not $\Theta(3^n)$.

$$2^n = o(3^n)$$

$$n^{100} < 2^n$$

$$\lim_{n \rightarrow \infty} \frac{(2^7)^{100} \cdot n^7}{2^7} = \infty$$

for $n = 2 \cdot 100$

$$n = \left(2^{14}\right)$$

$$\underbrace{(2^{100})^{100}}_{100} < \underbrace{(2^{100})^{100}}_{100}$$

$$n = (2^7)^{100}$$

$$(2^{700})^{100} < 2^{(2^{700})}$$

$$= 2^{70000}$$

$$\log n \quad \text{vs} \quad n^{0.01}$$

$$n = 10^{100}$$

$$100$$

$$(10^{100})^{0.01} = 10$$

1.29. Use the \prec relation to order the following functions by growth rate:

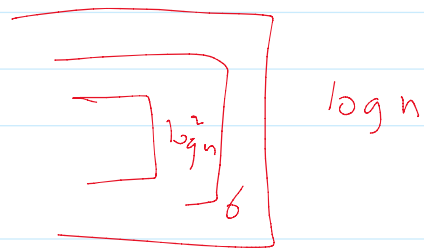
$$\underbrace{n^{1/100}}_{(6)}, \underbrace{\sqrt{n}}_{(7)}, \underbrace{\log n^{100}}_{(4)}, \underbrace{n \log n}_{(8)}, \underbrace{5}_{(2)}, \underbrace{\log \log n}_{(3)}, \underbrace{\log^2 n}_{(5)}, \underbrace{(\sqrt{n})^n}_{(9)}, \underbrace{(1/2)^n}_{(1)}, \underbrace{2^{n^2}}_{11}, \underbrace{n!}_{10}$$

$$(\sqrt{n})^n = (n^{1/2})^n = n^{n/2} = \underbrace{n \cdot n \cdot \dots \cdot n}_{n \text{ times}}$$

$$\underline{n!} = \underbrace{n \cdot (n-1) \cdot \dots \cdot (n/2) \cdot \dots \cdot 1}_{n \text{ terms}}$$

$$2^{(n^2)} \quad n^{\quad}$$

· $\text{ange}(1, n+1)$



$6 \log^3 n$

$$6 \times \frac{n(n+1)(2n+1)}{6}$$

2
4
8
16
32
64
128
256
512
1024

$2^{10} =$