

# Regular Languages

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## Recall: Regular Expressions

RE is a string on  $\Delta$

$\Delta = \Sigma \cup \{*, +, (, ), \emptyset\}$ , such that

1.  $\emptyset$  and  $\forall a \in \Sigma$  are the basic REs.
2. if  $\alpha$  is RE, then so is  $\alpha^*$  (star)
3. if  $\alpha$  and  $\beta$  are REs, then so is  $\alpha\beta$  (concatenation)
4. if  $\alpha$  and  $\beta$  are REs, then so is  $\alpha + \beta$  (union)

e.g.

$$L((ab)^*) = \{\lambda, ab, abab, ababab, \dots\}$$

$$L(ab^*) = \{a, ab, abb, abbb, \dots\}$$

2] Precedence rules:  $*$ , concatenation, union.

$$L(ab^* + b) = \{a, b, ab, abb, \dots\}$$

3] Notation: in Rule 4, the union operation '+' means "or", alternatively some textbooks use "|", "U"  
We will use + by default.

4] Notation:  $\alpha$  is a RE

$L(\alpha)$  denotes the language represented by  $\alpha$

e.g.

$$\alpha = a(ab)^*$$

$L(\alpha) =$  all strings start in a

$$\beta = (ab^*)^*$$

$L(\beta) =$  all strings start in a

$$L(\alpha) = L(\beta) \iff \alpha \equiv \beta$$

We can do  $\alpha = L(\alpha) = \beta$  if we talk about the languages

Exer:

Language L	$x \in L$	$y \notin L$	RE $\alpha$
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Exer:

Language L	$x \in L$	$y \notin L$	RE $\alpha$
① start in 1 end in 0 on $\{0,1\}$	10, 10010	101	$1(0+1)^*0$
② all strings on $\{a,b\}$	aa, bbb	none	$(a+b)^* \equiv (a^*b^*)^*$
③ $\{ \} = \emptyset$	none	ab, a, $\lambda$	$\emptyset$
④ $\{ \lambda \}$	$\lambda$	ab, a	$\emptyset^*$
⑤ $\{w \mid w \text{ has even a's}\}$	b, aa, aba baaaabb	abb,	$b^*(b^*ab^*ab^*)^*$ $\equiv b^*(ab^*ab^*)^*$
⑥ $\{w \mid w \text{ ends in b and has even b's}\}$	abaaab	abba, $\lambda$	$(a^*ba^*b)(a^*ba^*b)^*$

5] Def<sup>n</sup>. L is a regular language if it has a RE  $\alpha$  s.t.  $L(\alpha) = L$ .

e.g. see above.

e.g. Is  $L_p = \{w \mid w = w^p\}$  regular?

$L_p$  Not regular -

e.g.  $L = \{w \mid w = a^n b^n \text{ for } n \geq 1\}$

e.g. aaabbb

Not regular

$a^*b^*$

ab + aabbb +

aaabbb + ...

6] Thm: Regular languages are closed under

union,

$\left. \begin{array}{l} L_1 = L(\alpha_1) \quad L_2 = L(\alpha_2) \\ \alpha_1 + \alpha_2 \end{array} \right\}$

union,  
intersection,  
complement  
concatenation,  
and star operations

$L_1 = L(\alpha_1)$   $L_2 = L(\alpha_2)$   
 $\alpha_1 + \alpha_2$   
 $\alpha_1 \alpha_2$   
 $(\alpha_1)^*$

7] Thm: any finite language is regular?

$$L = \{w_1, w_2, w_3, \dots, w_n\}$$