

Objective: to introduce the topic of the theory of Computation, which is the mathematical study of computing machines and their capabilities.

## Lecture Notes on Theory of Computation (LNTC)

### §1. Alphabet and languages

language  $\rightarrow$  words  $\rightarrow$  Letters  
 each letter  $\in$  Alphabet  
 each word  $\in$  Language

e.g. English Alphabet (A) //  $A = \{a, b, c, \dots, z\}$

English language (L)

$a, p, l, e \in A$

$\downarrow$  make: "apple"  $\in L$

$\Rightarrow$  apple is a word in L

$\downarrow$  aaea  $\notin L \Rightarrow$  aaea not a word.

1] An alphabet is a finite set of symbols, denoted by  $\Sigma$

e.g.  $\Sigma = \{a, b, c\}$  // has 3 symbols

$\Sigma = \{0, 1\}$  // has 2 symbols.

$\Sigma = \{a, <, \$, \#, \text{!}, @, \&\}$  // anything can be a symbol.

$\Sigma = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, \dots\}$   $\leftarrow$  not valid.

$\Sigma = \{a\}$  // valid

$\Sigma = \{\} = \emptyset$  // valid, it could be empty;

2] A string on  $\Sigma$  is a finite sequence of symbols from  $\Sigma$ .

e.g.

$\Sigma = \{a, b, c\}$  : aabba, aa, aa, aa, aa, ...

e.g.

for  $\Sigma = \{a, b, c\}$  :  $abb, ac, cc$  are strings  
 $\Sigma = \{a\}$  :  $aa, aaaa, a$  are strings

3] The length of a string  $s$ ,

$|s|$  = number of symbols in  $s$ .

e.g.

$$s = 1100 \implies |s| = 4$$

$$s = aaa \implies |s| = 3$$

4] The empty string (or null) is a string of length zero  
• it is defined on every alphabet.  
• denoted by  $\lambda$

$$|\lambda| = 0$$

5] Note: any object can be a symbol. However, we use common characters like  $\{0, 1\}$ ,  $\{a, b\}$  ← our default.

6] Concatenation: of two strings  $x$  and  $y$  on some  $\Sigma$  is

$x \cdot y$  or simply  $xy$  (Juxtaposition)

e.g.  $x = abb$ ,  $y = aaa$

$$x \cdot y = abbaaa$$

$$yx = aaaabb$$

$$xx = abbaabb$$

$$x \cdot \lambda = x$$

$$\lambda \cdot x = x$$

7] Notation:  $s^n = \underbrace{s \cdot s \cdot s \cdot \dots \cdot s}_{n \text{ - times}}$

$$s^0 = \lambda$$

$$\text{Recursively, } s^n = \begin{cases} \lambda & \text{if } n=0 \\ s \cdot s^{n-1} & \text{if } n>0 \end{cases}$$

e.g.  $x = abb$

$$x^4 = abbabbabbabb$$

$$(aa)^3 = aaaaaa$$

$$a^3 = aaa$$

8] In general, let  $A$  be a set of symbols, then

$$A^n = \{s \mid s \text{ is a string on } A \text{ of length } n\}$$

e.g.  
 $A = \{b, c\}$

$$A^3 = \{bbb, bbc, bcb, bcc, cbb, cbc, ccb, ccc, \dots\}$$

e.g.  
 $\{a, b, c\}^2 = \{aa, ab, ac, ba, bb, bc, ca, cb, cc\}$

e.g.  $\Sigma^n$  is the set of all strings of length  $n$  on  $\Sigma$ .

9] Kleene star operation: let  $a \in \Sigma$ , then

$a^*$  is the set of all strings made of  $a$  repeated zero or more times

$$a^* = \{s \mid s = a^n \text{ for all } n \geq 0\}$$

$$= \{ \lambda, a, aa, aaa, \dots \}$$

10] In general the Kleene star operator can be applied on any set of symbols  $A$ ,

$$A^* = \{ s \mid s \text{ is a string on } A \}$$

e.g. ①  $\Sigma^*$  is the set of all possible strings on  $\Sigma$ .

$$\textcircled{2} \{a, c\}^* = \{ \lambda, a, c, aa, ac, ca, cc, aaa, \dots \}$$

Exer:

$$a^* \cup b^* = \{a, b\}^* \quad \text{No, for } aba \in \{a, b\}^* \text{ but not in LHS}$$

$$a^* \cap b^* = \emptyset \quad \text{False, for } \lambda \in \text{LHS but not in } \emptyset$$

11] Substrings:

string  $y$  is a substring of  $w$  iff  $\exists x, z \in \Sigma^*$  s.t.

$$w = xyz$$

$y$  is called prefix of  $w$  if  $x$  is null.

$y$  is called suffix of  $w$  if  $z$  is null.

e.g.  $aa$  is substring of  $baaab$  for  $x=b, z=ab$

$ab$  is suffix of  $baaab$  for  $x=baa, z=\lambda$

