

Chinese Remainder Theorem

Wednesday, February 9, 2022 12:46 PM

Recall: Linear Congruence
 $a \cdot x \equiv b \pmod{n}$

Missing:
#2, 5, 12, 22, 24
26, 32,

1] The Chinese Remainder Theorem (CRT)

Objective: to solve a system of linear congruences (in different mod's)

Thm: Let m_1, m_2, \dots, m_n be pairwise relatively prime integers. Then the system:

$$\begin{aligned}x &\equiv a_1 \pmod{m_1} \\x &\equiv a_2 \pmod{m_2} \\&\vdots \\x &\equiv a_n \pmod{m_n}\end{aligned}$$

has a unique solution modulo $m = m_1 \cdot m_2 \cdot \dots \cdot m_n$

2] Proof: How to find the solution in the CRT

① let $M_k = m / m_k \implies \gcd(M_k, m_k) = 1$

② $y_k = M_k^{-1} \pmod{m_k} \implies M_k \cdot y_k \equiv 1 \pmod{m_k}$

③ then $x \equiv a_1 M_1 y_1 + a_2 M_2 y_2 + \dots + a_n M_n y_n \pmod{m}$

3] e.g. on CRT (by Chinese mathematician Sun-Tsu)

Solve:

$$\begin{aligned}x &\equiv 2 \pmod{3} \\x &\equiv 3 \pmod{5} \\x &\equiv 2 \pmod{7}\end{aligned}$$

Solⁿ.

$$\begin{aligned}M_1 &= 5 \cdot 7 = 35 \implies y_1 \equiv 35^{-1} \equiv 2^{-1} \equiv 2 \pmod{3} \\M_2 &= 3 \cdot 7 = 21 \implies y_2 \equiv 21^{-1} \equiv 1^{-1} \equiv 1 \pmod{5} \\M_3 &= 3 \cdot 5 = 15 \implies y_3 \equiv 15^{-1} \equiv 1^{-1} \equiv 1 \pmod{7}\end{aligned}$$

$$\begin{aligned}x &\equiv 2 \cdot 35 \cdot 2 + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1 \pmod{105} \\&\equiv (3+1) \cdot 35 + (5-2) \cdot 21 + (7-5) \cdot 15\end{aligned}$$

$$\equiv 35 - 42 + 30 \equiv 23 \pmod{105}$$

4] Computer Arithmetic with large integer: (RNS)

Let m_1, m_2, \dots, m_n be pairwise relatively prime, then by CRT, $\forall a \in \mathbb{Z}_m$, $m = \prod m_i$, a can be represented uniquely by the n -tuple

$$(a \bmod m_1, a \bmod m_2, \dots, a \bmod m_n)$$

e.g. for $\mathbb{Z}_{12} = \{0, 1, \dots, 11\}$, use $m_1 = 3, m_2 = 4$

$0 = (0, 0)$	$4 = (1, 0)$	$8 = (2, 0)$
$1 = (1, 1)$	$5 = (2, 1)$	$9 = (0, 1)$
$2 = (2, 2)$	$6 = (0, 2)$	$10 = (1, 2)$
$3 = (0, 3)$	$7 = (1, 3)$	$11 = (2, 3)$

Add

$$\begin{array}{r} 6 \longrightarrow (0, 2) \\ 5 + \longrightarrow (2, 1) \\ \hline 11 \xleftarrow{\text{CRT}} (2, 3) \end{array}$$

Mult

$$\begin{array}{r} 2 \longrightarrow (2, 2) \\ 5 \times \longrightarrow (2, 1) \\ \hline 10 \xleftarrow{\text{CRT}} (1, 2) \end{array}$$

5] Exer. To perform arithmetic quickly on a CPU of max-int < 100 . Design an RNS for int $\approx 10^6$

Sol.

$$\text{Let } m_1 = 99, m_2 = 98, m_3 = 97, m_4 = 95 \Rightarrow m = 89 \times 10^6$$

$$123684 \longrightarrow (33, 8, 9, 89)$$

$$413456 + \longrightarrow (32, 92, 2, 16)$$

$$\boxed{\cdot x}$$

$$(65, 2, 11, 10)$$

Solve

by CRT

Solve

$$\begin{cases} x \equiv 65 \pmod{99} \\ x \equiv 2 \pmod{98} \\ \vdots \\ x \equiv 10 \pmod{95} \end{cases} \left. \vphantom{\begin{cases} x \equiv 65 \pmod{99} \\ x \equiv 2 \pmod{98} \\ \vdots \\ x \equiv 10 \pmod{95} \end{cases}} \right\} \text{ by CRT}$$

6] Hash functions :

$h(x) \rightarrow y$ where x can be of any length
and y has a fixed length

e.g.

$$h(x) = x \pmod{31}$$

the h output is of length 5 bits

e.g. $h(65) = 3 \rightarrow (00011)_2$

7] Pseudorandom numbers

modulus m
multiplier a , $2 \leq a < m$
increment c , $0 \leq c < m$
Seed $x_0 \in \mathbb{Z}_m$

$$x_i = (a x_{i-1} + c) \pmod{m}$$

e.g. $m=9$, $a=7$, $c=4$, $x_0=3$

$$x_1 = 7 \cdot (3) + 4 \equiv 7 \pmod{9}$$

$$x_2 = 7 \cdot 7 + 4 \equiv 8 \pmod{9}$$

\vdots